



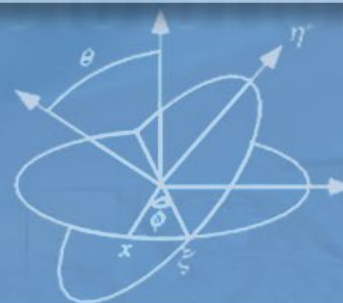
JHU vision lab

Scalable Subspace Clustering

Chong You

Johns Hopkins University

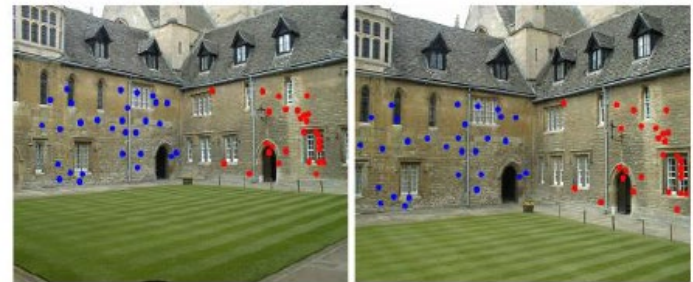
Joint work with Chun-guang Li, Daniel P. Robinson, and René Vidal



Motivation

In many areas, we deal with large amount of data

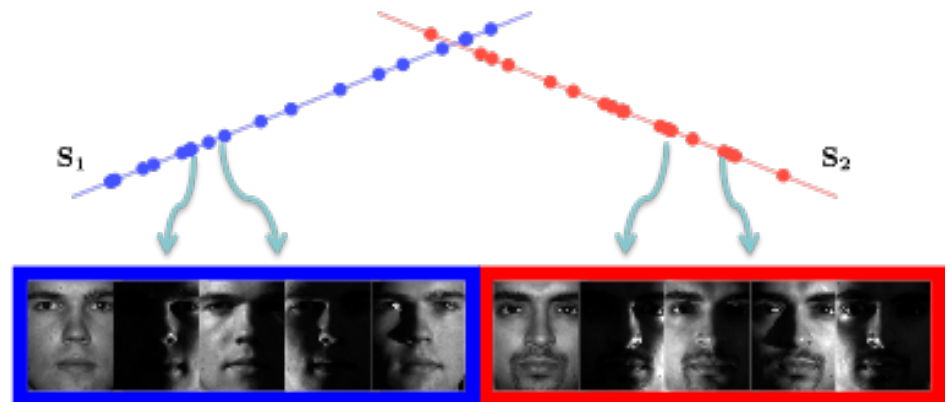
- Data contains **multiple classes**
- Each class lies in a **low-dimensional subspaces**



Planar Segmentation



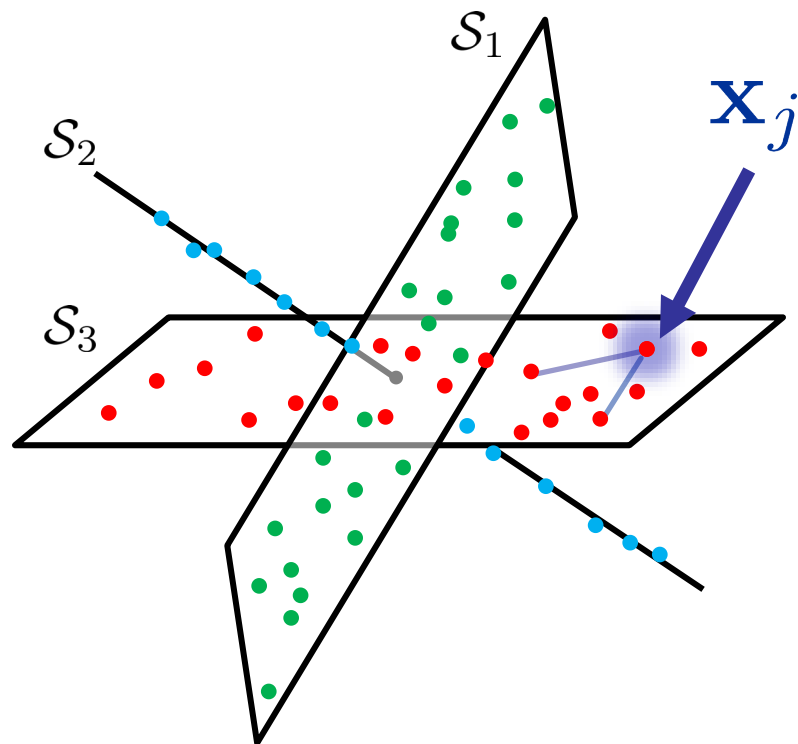
Motion Segmentation



Face Recognition/Clustering

Subspace clustering

Given data $X = [\mathbf{x}_1, \dots, \mathbf{x}_N]$, find a union of subspaces that fits the data:



Two-step Approach

- Build data affinity
- Apply spectral clustering

Challenges

- Distance based affinity fails at the intersection of subspaces

Self-Expressive Model

- Compute affinity by data self-representation

Prior work: sparse subspace clustering (SSC)

Sparsity

$$\min_{\mathbf{c}_j} \|\mathbf{c}_j\|_0 \quad \text{s.t.}$$

Self-representation

$$\mathbf{x}_j = X\mathbf{c}_j, c_{jj} = 0$$

Convex relaxation



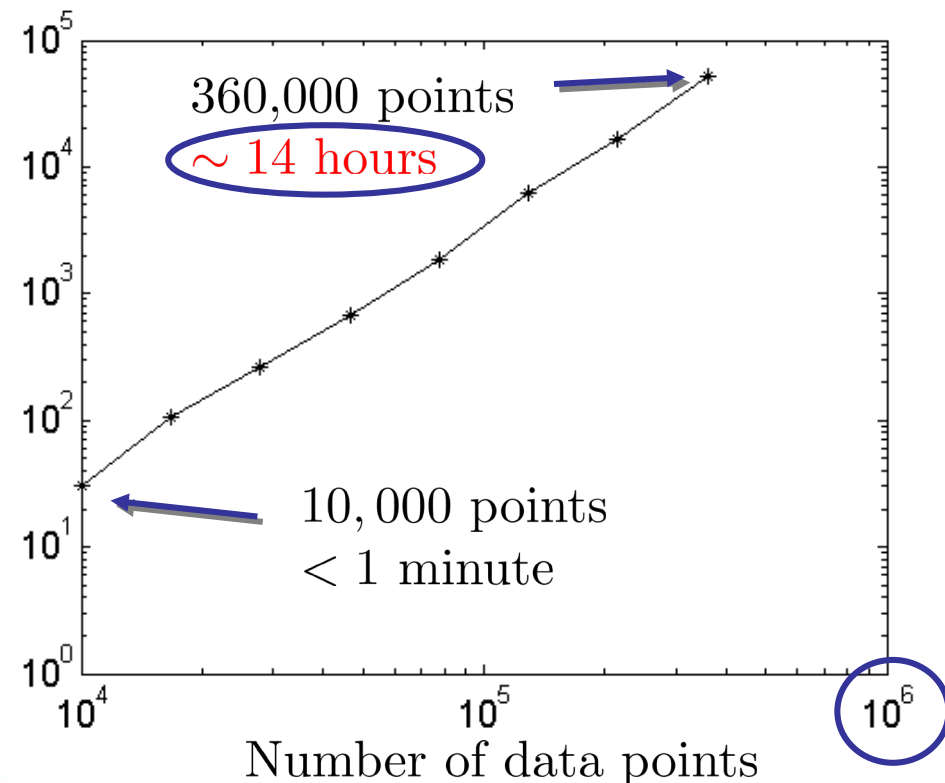
Method:

SSC by basis pursuit
(SSC-BP)

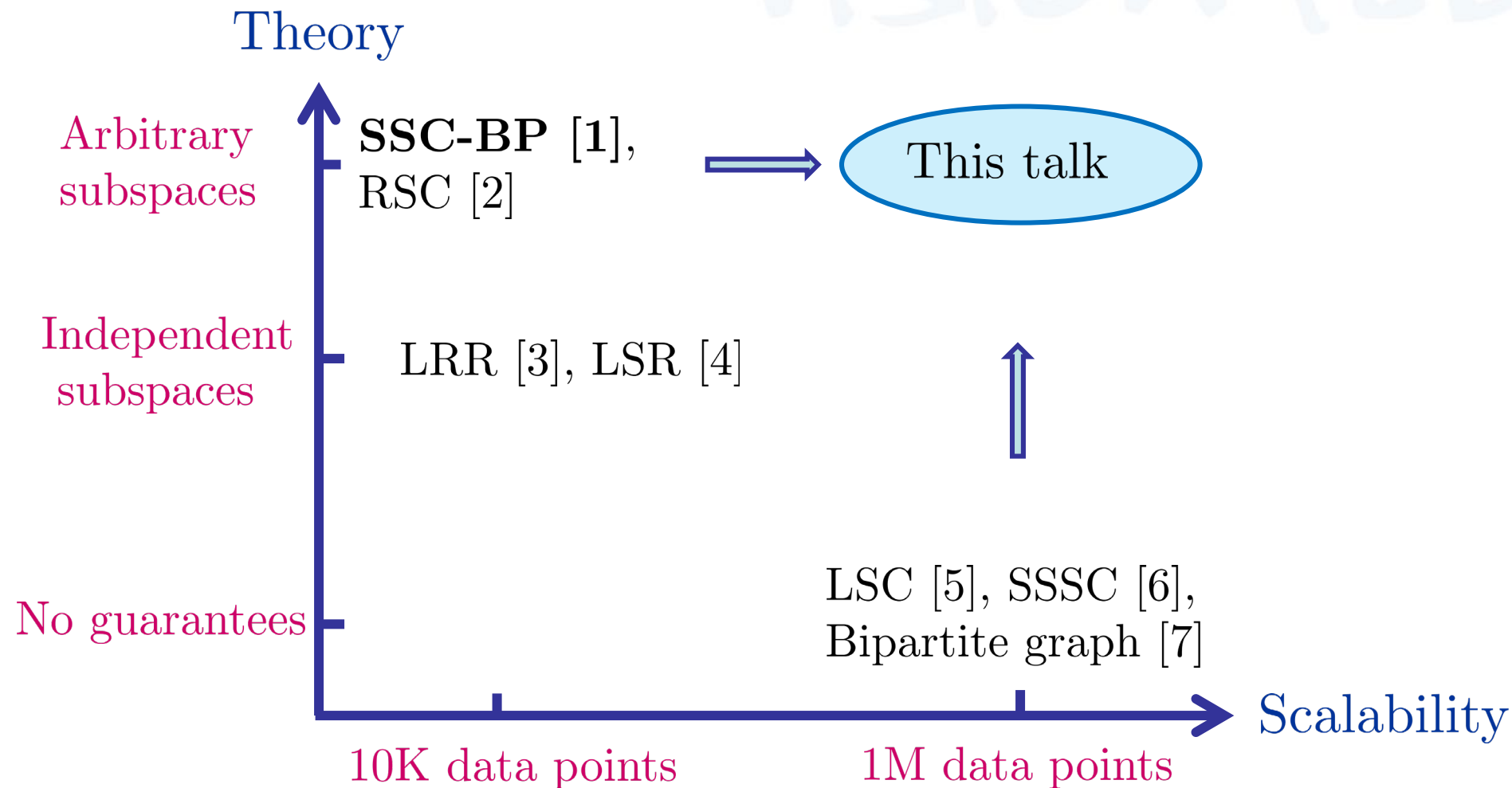
Properties:

- ✓ Guaranteed correct connections
- ✗ Not scalable:
solved by CVX/ADMM
tested on ≤ 640 points

Running time (sec.)



Prior work: overview



- [1] E. Elhamifar and R. Vidal, Sparse Subspace Clustering, CVPR'09
- [2] M. Soltanolkotabi and E. Candes, Robust Subspace Clustering, Annual of Statistics'13
- [3] G. Liu, Z. Lin, Y. Yu, Robust Subspace Segmentation by Low-Rank Representation, ICML'10
- [4] Lu et al., Robust and efficient subspace segmentation via least squares regression, ECCV 2012.
- [5] X. Chen and D. Cai, Large Scale Spectral Clustering with Landmark-based Representation, AAAI'11
- [6] X. Peng, L. Zhang, Z. Yi, Scalable Sparse Subspace Clustering, CVPR'13
- [7] A. Adler, M. Elad, Y. Hel-Or, Linear-Time Subspace Clustering via Bipartite Graph Modeling



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Scalable Sparse Subspace Clustering by Orthogonal Matching Pursuit

Chong You[†], Daniel P. Robinson[‡], René Vidal[†]

[†]Center for Imaging Science, Johns Hopkins University

[‡]Applied Mathematics and Statistics, Johns Hopkins University



Sparse subspace clustering

Sparsity

$$\min_{\mathbf{c}_j} \|\mathbf{c}_j\|_0 \quad \text{s.t.}$$

Self-representation

$$\mathbf{x}_j = X\mathbf{c}_j, c_{jj} = 0$$

Convex relaxation



Method:

SSC by *basis pursuit*
(SSC-BP)

Properties:

- ✓ Guaranteed correct connections
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SSC by Orthogonal Matching Pursuit

Sparsity

$$\min_{\mathbf{c}_j} \|\mathbf{c}_j\|_0$$

s.t.

Self-representation

$$\mathbf{x}_j = X\mathbf{c}_j, c_{jj} = 0$$

Convex relaxation



Method:

SSC by *basis pursuit*
(SSC-BP)

Properties:

- ✓ Guaranteed correct connections
- ✗ Not scalable:
solved by CVX/ADMM
tested on ≤ 640 points

Greedy pursuit



Method:

SSC by *orthogonal matching pursuit*
(SSC-OMP)

Properties:

- ? Guaranteed correct connections
- ? Scalable

SSC by Orthogonal Matching Pursuit

Sparsity

$$\min_{\mathbf{c}_j} \|\mathbf{c}_j\|_0$$

s.t.

Self-representation

$$\mathbf{x}_j = X\mathbf{c}_j, c_{jj} = 0$$

Convex relaxation



Method:

SSC by *basis pursuit*
(SSC-BP)

Properties:

- ✓ Guaranteed correct connections
- ✗ Not scalable:
solved by CVX/ADMM
tested on ≤ 640 points

Greedy pursuit



Method:

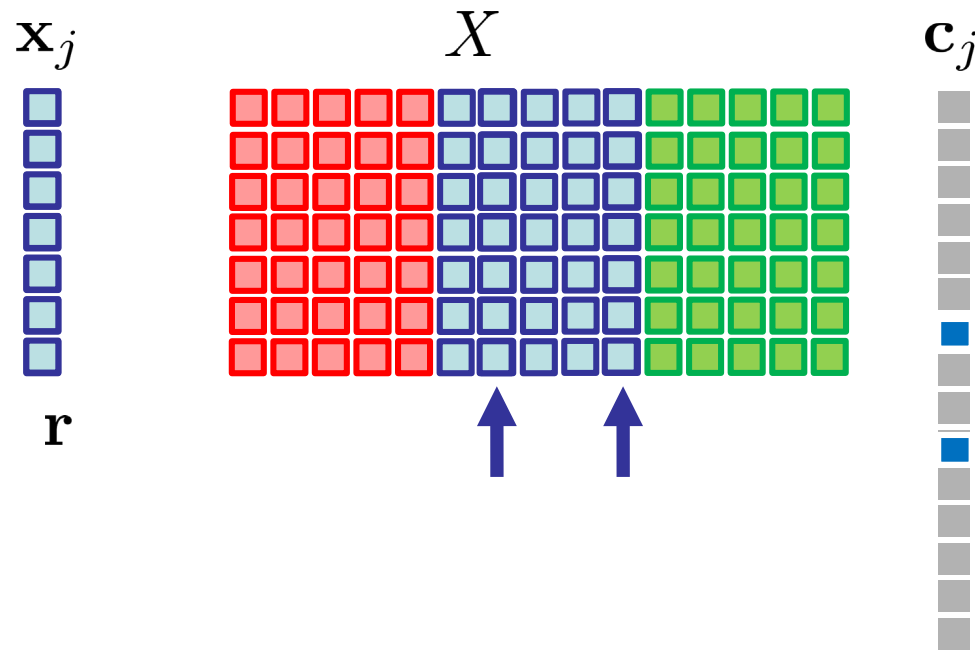
SSC by *orthogonal matching pursuit*
(SSC-OMP)

Contributions:

- ✓ Guaranteed correct connections
- ✓ Scalable:
tested on 1,000,000 points

SSC by Orthogonal Matching Pursuit

Find representation $\mathbf{x}_j = X\mathbf{c}_j$ by greedy selection



What are the conditions for giving correct connections?
Each iteration picks a point from the same subspace

Guaranteed correct connections: deterministic model

Theorem

Suppose that $\mathbf{x}_j \in \mathcal{S}_\ell$. Then, \mathbf{c}_j gives correct connections if

$$\mu(W_j^\ell, X^{-\ell}) < r^\ell,$$

where μ captures the **similarity** between \mathcal{S}_ℓ and all other subspaces, and r captures **distribution** of points in \mathcal{S}_ℓ .

For SSC-BP^[3]:

$W_j^\ell =$ dual points

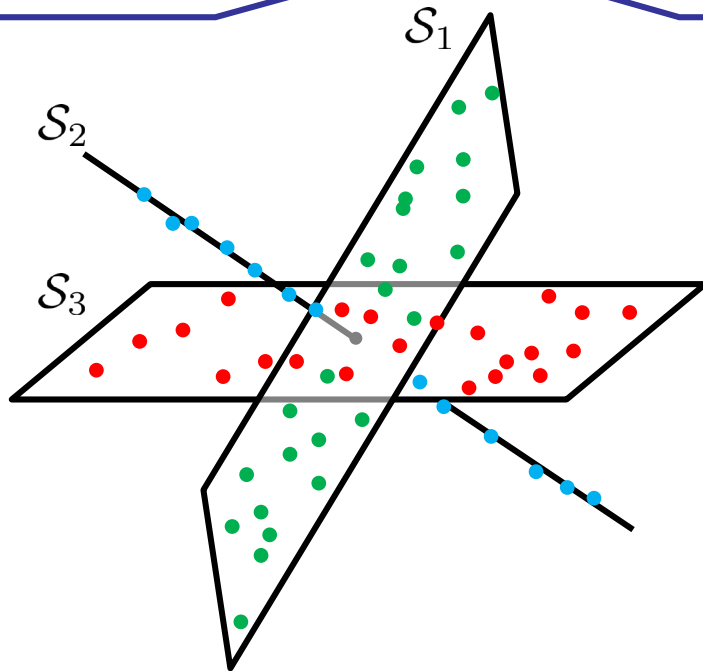
For SSC-OMP:

$W_j^\ell =$ residual points

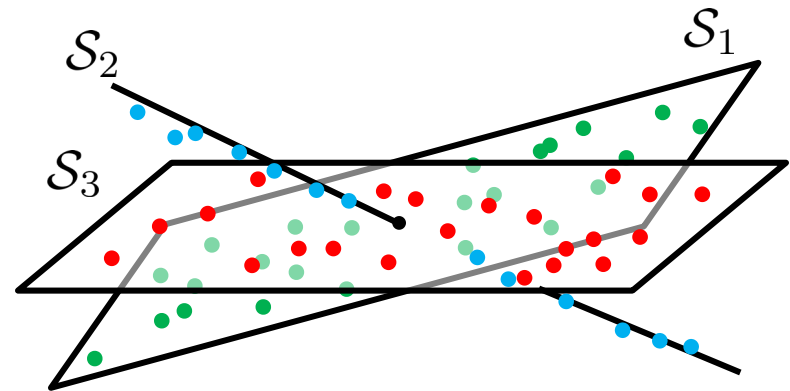
Guaranteed correct connections: deterministic model

$$\mu(W_j^\ell, X^{-\ell}) < r^\ell$$

Similarity between subspaces



Easier case



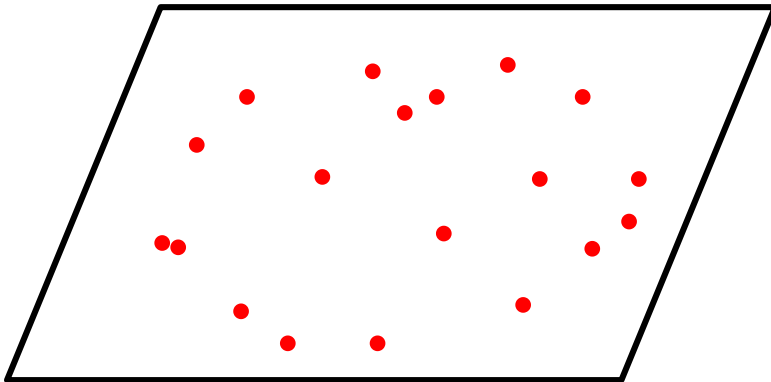
Harder case

Guaranteed correct connections: deterministic model

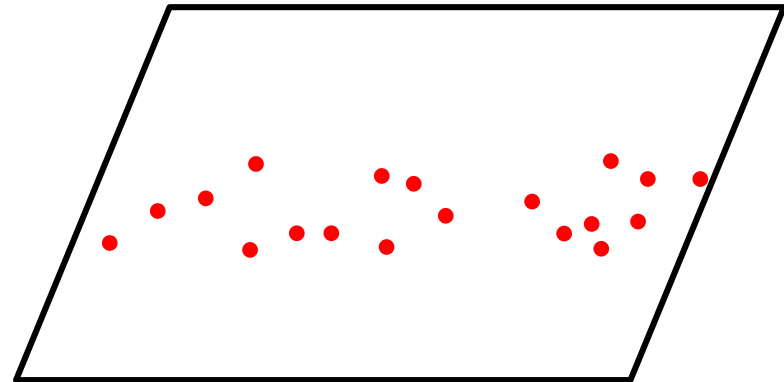
$$\mu(W_j^\ell, X^{-\ell}) < r^\ell$$

Similarity between subspaces

Distribution of points



Easier case



Harder case

Is this condition likely to be satisfied ?

Guaranteed correct connections: random model

Random model:

- Draw n subspaces of dimension d in ambient dimension D
- Draw $\rho d + 1$ points from each subspace

Theorem

Under the random model, the solution $\{\mathbf{c}_j\}_{j=1}^N$ gives correct connections with overwhelming probability if

$$\frac{d}{D} < \frac{c^2(\rho) \log \rho}{12 \log N}$$

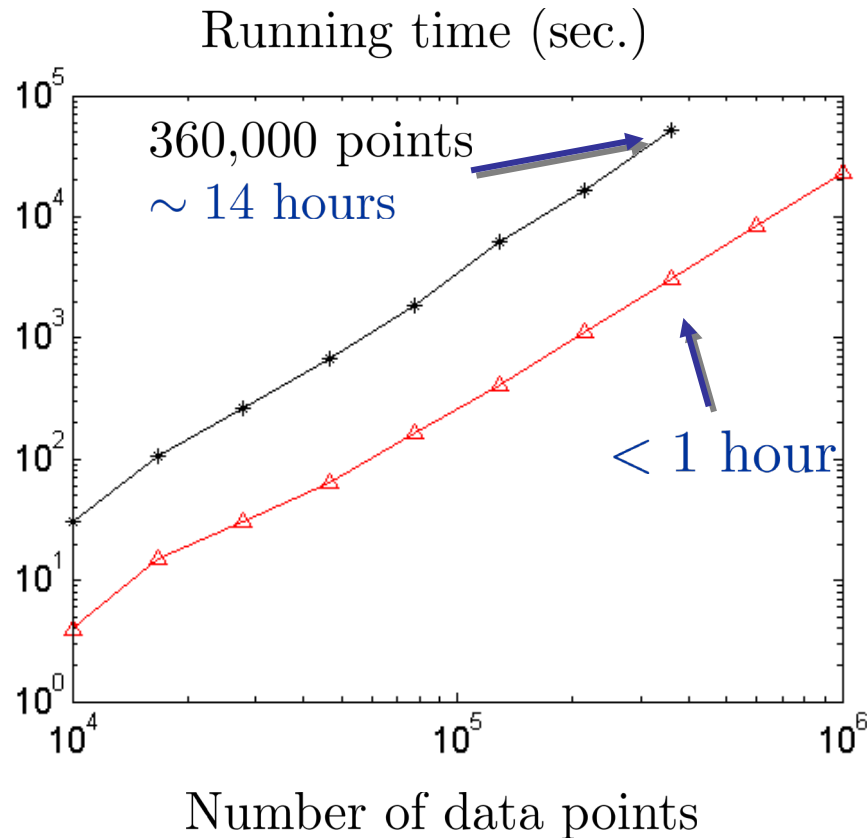
For SSC-BP^[3]:

$$p > 1 - \frac{2}{N} - Ne^{-\sqrt{\rho}d}$$

For SSC-OMP:

$$p > 1 - \frac{2d}{N} - Ne^{-\sqrt{\rho}d}$$

Scalability: SSC-BP versus SSC-OMP



- SSC-BP (Baseline)
- SSC-OMP

SSC-OMP significantly reduces the time, and deals with 1 million data

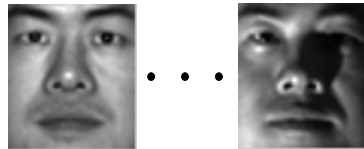
Experiment on extended Yale B

img-1 ... img-64

subject-1



subject-2



...

subject-38



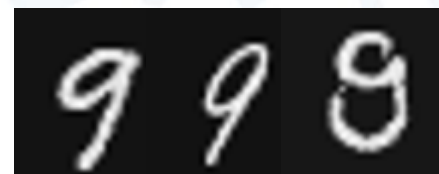
No. subjects	2	10	20	30	38
<i>a%: average clustering accuracy</i>					
SSC-OMP	99.21	88.43	81.71	79.27	80.45
SSC-BP	99.45	91.85	79.80	76.10	68.97
LSR	96.77	62.89	67.17	67.79	63.96
LRSC	94.32	66.98	66.34	67.49	66.78
SCC	78.91	NA	NA	14.15	12.80
<i>t(sec.): running time</i>					
SSC-OMP	0.3	1.7	4.7	9.4	14.5
SSC-BP	49.1	228.2	554.6	1240	1851
LSR	0.1	0.8	3.1	8.3	15.9
LRSC	1.1	1.9	6.3	14.8	26.5
SCC	50.0	NA	NA	520.3	750.7

> 100 times faster

Experiment on MNIST



...



No. points	500	2,000	6,000	20,000	60,000
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a%: average clustering accuracy

SSC-OMP	85.17	88.99	90.56	94.21	94.68
SSC-BP	83.01	85.58	85.60	-	-
LSR	75.84	78.09	79.91	-	-
LRSC	75.02	79.44	79.88	-	-
SCC	53.45	66.43	70.60	-	-

t(sec.): running time

SSC-OMP	1.3	11.7	71.7	427	3219
SSC-BP	20.1	635.2	13605	-	-
LSR	1.7	42.4	327.6	-	-
LRSC	1.9	43.0	312.9	-	-
SCC	31.2	101.3	366.8	-	-

SSC by Orthogonal Matching Pursuit (OMP):



theoretical guarantees for **correct connections**



performance validation on **large databases**



JHU vision lab

Scalable Elastic Net Subspace Clustering

Chong You[†], Chun-Guang Li^{*}, Daniel P. Robinson[‡], René Vidal[†]

[†]Center for Imaging Science, Johns Hopkins University

^{*}SICE, Beijing University of Posts and Telecommunications

[‡]Applied Mathematics and Statistics, Johns Hopkins University



Motivation

SSC

$$\min_{\mathbf{c}} \|\mathbf{c}\|_1 + \frac{\gamma}{2} \|\mathbf{x} - X\mathbf{c}\|_2^2$$

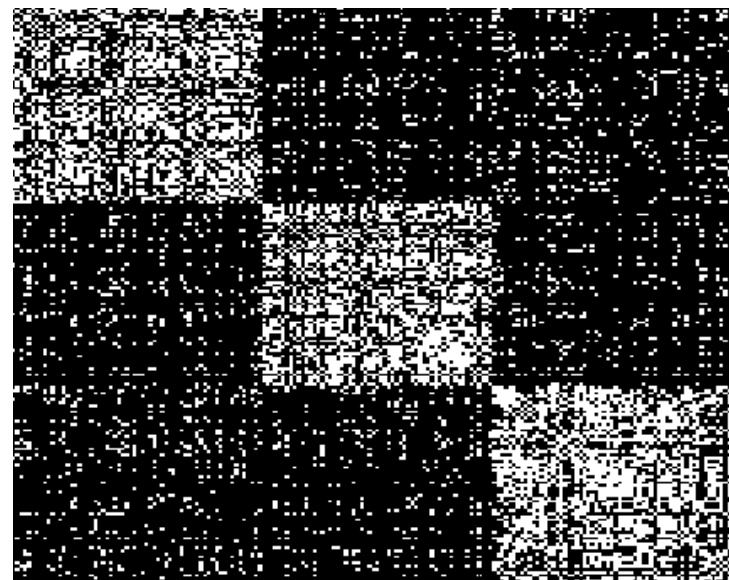


✓ Few wrong connections

✗ Not well connected

LSR

$$\min_{\mathbf{c}} \|\mathbf{c}\|_2^2 + \frac{\gamma}{2} \|\mathbf{x} - X\mathbf{c}\|_2^2$$

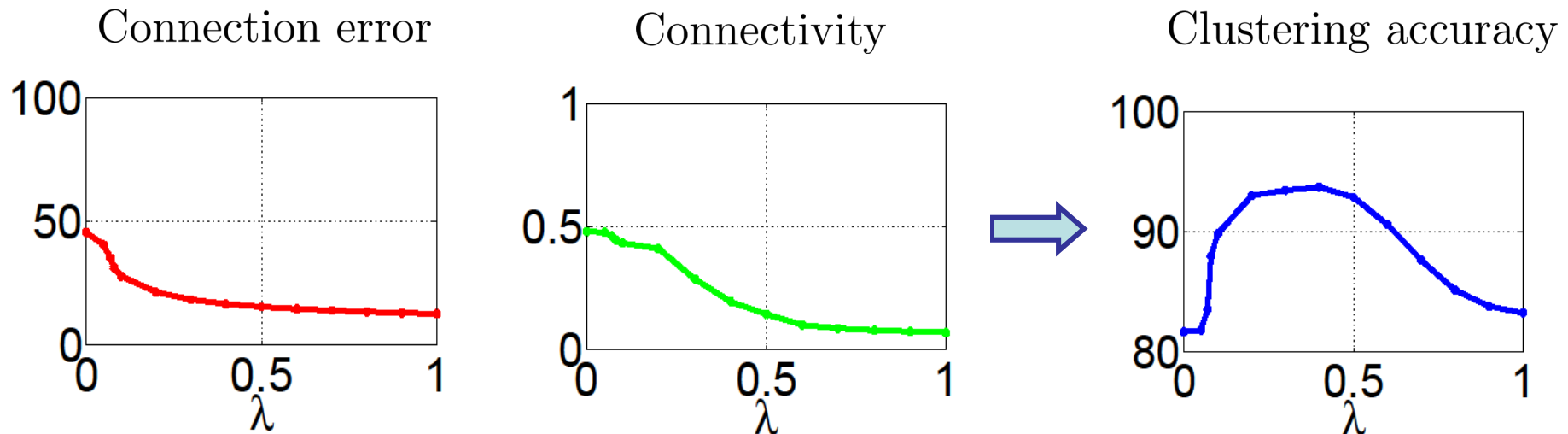


✗ Many wrong connections

✓ Well-connected

Elastic net Subspace Clustering (EnSC)

$$\min_{\mathbf{c}_j} \lambda \|\mathbf{c}_j\|_1 + \frac{1-\lambda}{2} \|\mathbf{c}_j\|_2^2 + \frac{\gamma}{2} \|\mathbf{x}_j - X\mathbf{c}_j\|_2^2 \quad \text{s.t.} \quad \mathbf{c}_{jj} = 0$$



Key theoretical challenges:

- ? Is EnSC guaranteed to give **correct connections**
- ? How to explain the tradeoff with **connectivity**

Scalable Elastic net Subspace Clustering

$$\min_{\mathbf{c}_j} \lambda \|\mathbf{c}_j\|_1 + \frac{1-\lambda}{2} \|\mathbf{c}_j\|_2^2 + \frac{\gamma}{2} \|\mathbf{x}_j - X\mathbf{c}_j\|_2^2 \quad \text{s.t.} \quad \mathbf{c}_{jj} = 0$$

- Prior methods
 - ADMM
 - Interior point
 - Solution path
 - Proximal gradient method
 - etc.

Scalability issue:

- Too many iterations to converge
- Access to full data matrix

Key challenge:

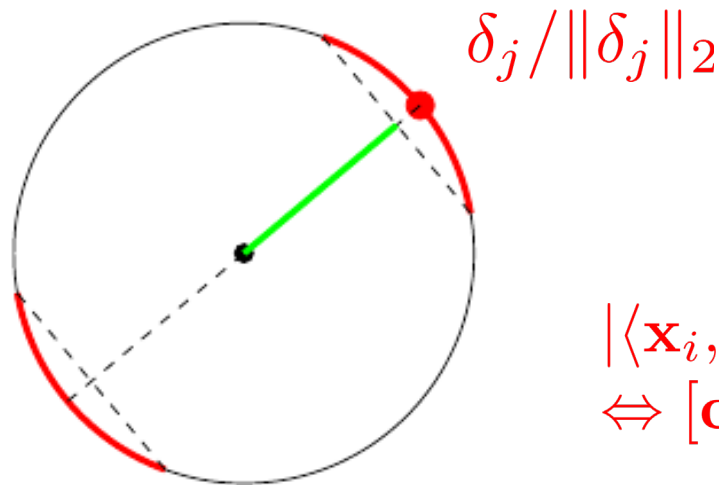
? Can we derive **scalable algorithms** that can handle 1 million data

Geometry of solution

$$\min_{\mathbf{c}_j} \lambda \|\mathbf{c}_j\|_1 + \frac{1-\lambda}{2} \|\mathbf{c}_j\|_2^2 + \frac{\gamma}{2} \|\mathbf{x}_j - X\mathbf{c}_j\|_2^2 \quad \text{s.t.} \quad \mathbf{c}_{jj} = 0$$

Oracle point $\delta_j = \gamma(\mathbf{x}_j - X\mathbf{c}_j^*)$

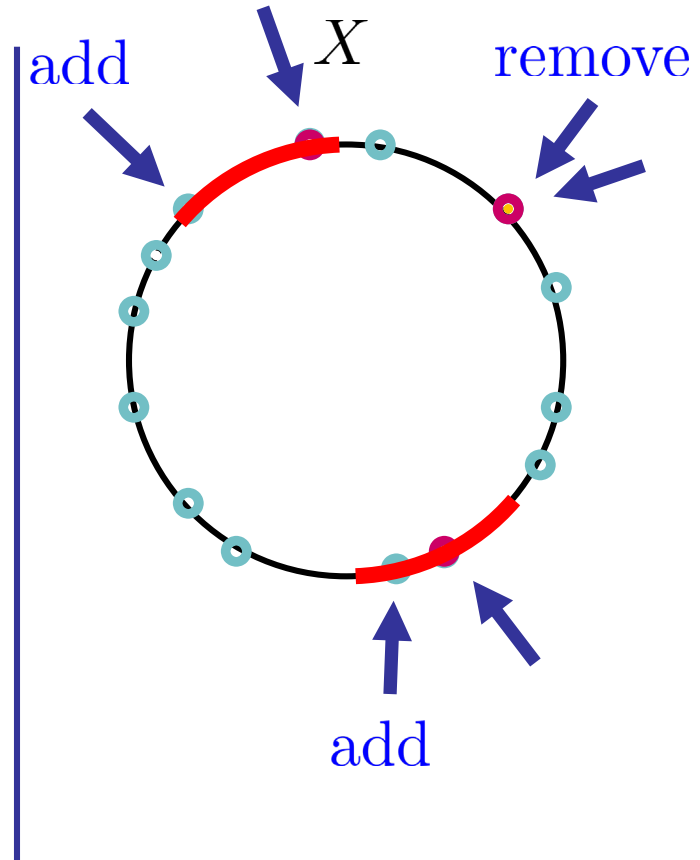
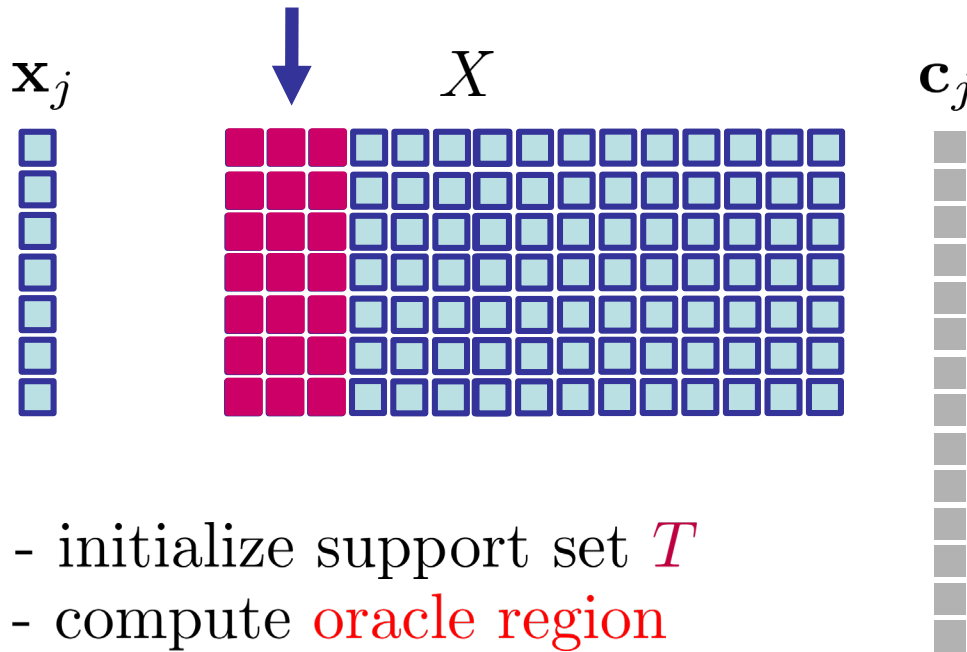
- If we know the solution \mathbf{c}_j^* , we can compute δ_j
- If we know δ_j , we can find the support of the solution \mathbf{c}_j^*



$$|\langle \mathbf{x}_i, \delta_j \rangle| > \lambda \\ \Leftrightarrow [\mathbf{c}_j^*]_i \neq 0$$

Oracle guided active set (ORGEN) algorithm

$$\min_{\mathbf{c}_j} \lambda \|\mathbf{c}_j\|_1 + \frac{1-\lambda}{2} \|\mathbf{c}_j\|_2^2 + \frac{\gamma}{2} \|\mathbf{x}_j - X\mathbf{c}_j\|_2^2 \quad \text{s.t.} \quad \mathbf{c}_{jj} = 0$$



Oracle guided active set (ORGEN) algorithm

$$\min_{\mathbf{c}_j} \lambda \|\mathbf{c}_j\|_1 + \frac{1-\lambda}{2} \|\mathbf{c}_j\|_2^2 + \frac{\gamma}{2} \|\mathbf{x}_j - X\mathbf{c}_j\|_2^2 \quad \text{s.t.} \quad \mathbf{c}_{jj} = 0$$

\mathbf{x}_j



Theorem:

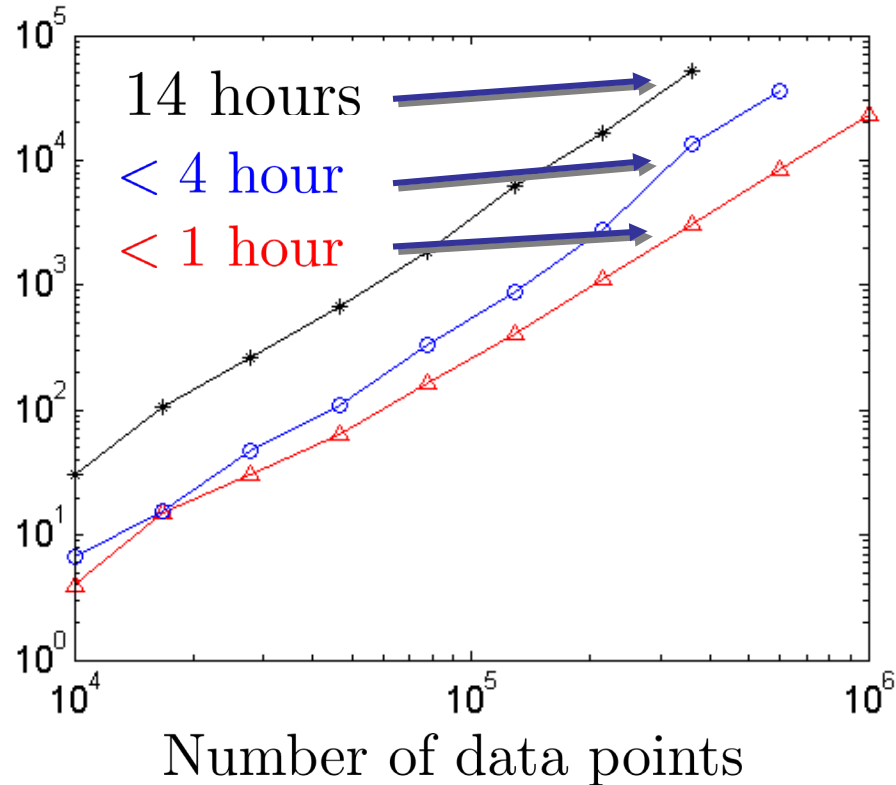
The support set T converges to the true support set in a finite number of iterations

- init
- con
- upc
- rep

Efficiency is gained by solving multiple **small** problems instead of one **big** problem

SSC-BP vs. SSC-OMP vs. EnSC-Oracle

Running time (sec.)



- SSC-BP (Baseline)
- SSC-OMP
- EnSC-Oracle

reduces the time for SSC-BP

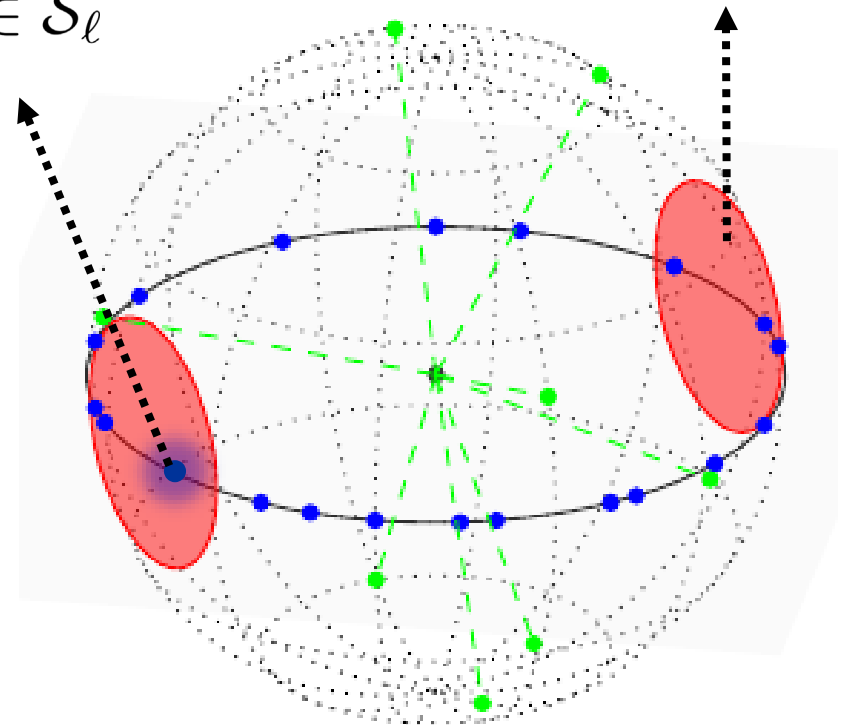
Correct connections vs. connectivity

$$\min_{\mathbf{c}_j} \lambda \|\mathbf{c}_j\|_1 + \frac{1-\lambda}{2} \|\mathbf{c}_j\|_2^2 + \frac{\gamma}{2} \|\mathbf{x}_j - X\mathbf{c}_j\|_2^2 \quad \text{s.t.} \quad \mathbf{c}_{jj} = 0$$

oracle region

- λ is large
 - \implies oracle region is small
 - \implies correct connection
- λ is small
 - \implies oracle region is large
 - \implies well-connected

$\mathbf{x}_j \in \mathcal{S}_\ell$



Guaranteed no wrong connections

$$\min_{\mathbf{c}_j} \|\mathbf{c}_j\|_1 + \frac{\gamma}{2} \|\mathbf{x}_j - X\mathbf{c}_j\|_2^2 \quad \text{s.t.} \quad \mathbf{c}_{jj} = 0$$

Theorem: (for SSC)

Condition for guaranteed no wrong connections:

$$\mu(W^{(\ell)}, X^{(-\ell)}) < r^{(\ell)}$$



Similarity between subspaces



Distribution of points

Guaranteed no wrong connections

$$\min_{\mathbf{c}_j} \lambda \|\mathbf{c}_j\|_1 + \frac{1-\lambda}{2} \|\mathbf{c}_j\|_2^2 + \frac{\gamma}{2} \|\mathbf{x}_j - X\mathbf{c}_j\|_2^2 \quad \text{s.t.} \quad \mathbf{c}_{jj} = 0$$

Theorem: (for EnSC)

Condition for guaranteed no wrong connections:

$$\mu(W^{(\ell)}, X^{(-\ell)}) < r^{(\ell)} - \frac{1-\lambda}{\lambda}$$



Similarity between subspaces

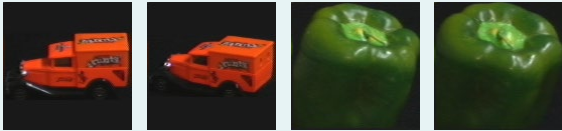




Distribution of points

Condition is **harder** to be satisfied
Graph has **better connectivity** } Higher clustering accuracy

Experiments

Test of EnSC with ORGEN on real data

database	# data	ambient dim.	# clusters	Examples
Coil-100	7,200	1024	100	
PIE	11,554	1024	68	
MNIST	70,000	500	10	
CovType	581,012	54	7	

Experiments

Our method (EnSC) achieves the best **clustering accuracy**

database	# data	SSC-BP	SSC-OMP	EnSC
Coil-100	7,200	57.10%	42.93%	69.24%
PIE	11,554	41.94%	24.06%	52.98%
MNIST	70,000	-	93.07%	93.79%
CovType	581,012	-	48.76%	53.52%

Experiments

Our method (EnSC) is **scalable**

database	# data	SSC-BP	SSC-OMP	EnSC
Coil-100	7,200	127 mins	3 mins	3 mins
PIE	11,554	412 mins	5 mins	13 mins
MNIST	70,000	-	6 mins	28 mins
CovType	581,012	-	783 mins	1452 mins

Conclusion

$$\min_{\mathbf{c}_j} \lambda \|\mathbf{c}_j\|_1 + \frac{1-\lambda}{2} \|\mathbf{c}_j\|_2^2 + \frac{\gamma}{2} \|\mathbf{x}_j - X\mathbf{c}_j\|_2^2 \quad \text{s.t.} \quad \mathbf{c}_{jj} = 0$$

- ✓ guaranteed correct connections
 - ✓ improved connectivity
 - ✓ efficient algorithm for large scale problems
- } better clustering

Acknowledgement

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Vision Lab @ Johns Hopkins University
<http://www.vision.jhu.edu>

Thank you!